



Exception Handling

Interfaces, Implementations, and Evaluation

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What do we want?

We want to produce programs which can

- detect exceptional conditions and
- *react to them.*

We also want these programs to be

- supported by our friendly neighborhood programming environments and
- amenable to optimization on current and future platforms.

- The standard specifies *interface requirements*
- Interfaces have multiple good implementations
- Good design is hard, but interface criteria include
 - Minimality
 - Orthogonality
 - Clarity
- We prefer explicit and local control and data flow

Outline

- **Motivating and rejected examples**
 - Deferring debugging...
- Survey of software interfaces
- Hardware support
- Hardware / software mapping
- A scarecrow proposal

Motivating / Archetypal Examples

- Algorithms that exceptions make risky
 - abort, wasting minimal work (Eureka exit)
 - then possibly do something else (complex multiply, scaling)
- Slightly change the arithmetic
 - substitute a limit for an exceptional result (continued fractions, replacement)
- Soften the arithmetic's boundaries
 - extend the dynamic range (long products, counting mode)
- Communicate the quality of a result

Rejected Examples

Any of the following could overly constrain our choices:

- Supporting heavy modifications to the arithmetic.
 - UN, OV, *etc.*
- Allowing extremely non-local, implicit control and data flow.
- Considering any particular hardware implementation.
- Requiring specific debugging tools...

Deferring Debugging

- General purpose tools handle uninitialized data.
 - Purify, valgrind, *etc.*
- Different applications need different retrospective diagnostic facilities.
- We're not sure how to support future debugging tools. (path-based, *etc.*)

We need to keep debugging in mind, but it is a “quality of implementation” issue.

Survey of Software Interfaces

- Motivating and rejected examples
- **Survey of software interfaces**
 - Try-catch
 - Flag testing
 - Explicit trapping
 - Substitution
 - Flag-carrying types
 - Conditional branching FP ops
- Hardware support
- Hardware / software mapping
- A scarecrow proposal

Try-Catch

```
try {  
    fast and sloppy code  
}  
catch (exceptional cases) {  
    slow and careful  
}
```

Floating-point mechanism exists in:

- fpmenu
- Ada
- Numerical Turing
- BASIC
- Common Lisp (optional)
- Borneo (specification)

Try-Catch

Language aspects:

- Scope is always specified as linguistic blocks.
- Extent:
 - Can called functions also raise exceptions?
 - Are “thrown” exceptions specified statically?
- How do callers / callees communicate which exceptions are interesting?
- Is the `try` block interrupted precisely?
- Can execution be resumed or statements restarted?

Try-Catch

Benefits:

- Matches existing, non-FP practice.
- Limits optimization impact to blocks.

Drawbacks:

- Existing practice is often mis-managed.

Try-Catch

Observations:

- Case without resumption / restart can be implemented through either traps or flags.
- Catching invalid is often followed by testing in-scope variables to determine *which* invalid op occurred.

Flag-Testing

```
double f (double x) {  
    save environment  
    do work;  
    if (flags raised) do alternate work;  
    restore environment  
    merge proper flags  
    return result  
}
```

Exists in:

- C99
- Many platform-dependent libraries
- Borneo (specification)

Flag-Testing

```
double f (double x) {  
    fenv_t fenv;  
    feholdexcept (&fenv);  
    do work;  
    if (flags raised) do alternate work;  
    fesetenv (&fenv);  
    return out1+out2;  
}
```

Exists in:

- C99
- Many platform-dependent libraries
- Borneo (specification)

Flag-Testing

Language aspects:

- Scope: Are flags set by block, or through a global datum?
- Extent: How do flags pass through subroutines?

Benefits:

- Predictable control flow.

Drawbacks:

- All operations share state.
- Subexp movement and compile-time evaluation often incorrect.
- Flag tests clutter code.

Flag-Testing

```
double f (double x) {  
    fenv_t fenv; feholdexcept (&fenv) ;  
    do work;  
    if (flags raised) do alternate work;  
    fesetenv (&fenv) ;  
    return out1+out2;  
}
```

Observations:

- Almost all uses follow the above pattern, including a few operations to set output flags implicitly.
- Compilers must virtualize and track flags for optimization.

Explicit Trapping

Implementations

- Sun's libm9x
- SIGFPE handling
(wmexcp, fpmenu)

Aspects

- Scope: dynamic
- Extent: dynamic

Benefits

- Unknown

Drawbacks

- No portable interfaces
- Nigh-impossible to use
- Serious non-local, implicit effects

Substitution

```
FPE_PRESUB (FE_INVALID, +INFINITY)
    for (i = 0; i < n_items; ++i)
        newprice[i] = price[i] + bidincr[i];
FPE_END_PRESUB
```

Exists in:

- IEEE defaults
- fpmenu: presub and counting

Substitution

```
FPE_COUNT (&cnt)
    for (i = 0; i < N; ++i)
        out *= A[i] + B[i];
FPE_END_COUNT
```

Exists in:

- IEEE defaults
- fpmenu: presub and counting

Substitution

Language aspects:

- Static scope, but static or dynamic extent
- How do you determine the replaced type?
- Do you consider operands? Get the sign?
- Location of count or other implicit operands?

Benefits:

- Well-defined, can have very limited scope
- Many implementation / optimization options

Drawbacks:

- Only two functionalities out of how many?

'New': Flag-Carrying Types

In the continued fraction code:

```
double f, f1, ...;
flagdouble r;
int j;
...
    r = d1/d;
    f1 = -r * d;
    if (!flagtest_and_clear(r, INVALID))
        continue;
    // fixup
...
```

- Explicit syntax for the desired result.
- Useful when only a few items are flagged.

Flag-Carrying Types

Language aspects:

- Scope and extent match value types'.
- Static typing = static flags
- Relies on expression evaluation typing

Flag-Carrying Types

Benefits;

- Everything is explicit.
- Optimizations use existing frameworks.
- User control over which expressions require flags.
- Programmers understand data-flow.

Drawbacks:

- Verbose (sub- and dynamic typing help)

Observations:

- Flagged compile-time constants keep flags.
- Subexpressions can be lifted.

'New': Conditional Branch FP-Ops

```
complex operator* (complex x, complex y)
{
    let
        double operator*(double, double) =
            trapping_mult(double, OVERFLOW: ov_label,
                          UNDERFLOW: un_label,
                          INVALID: not_complex_label);
        double operator+(double, double) =
            trapping_add(double, INVALID: infs_label);
        double operator-(double, double) =
            trapping_sub(double, INVALID: infs_label);
    in {
        return complex (real(x)*real(y) - imag(x)*imag(y),
                       real(x)*imag(y) - imag(x)*real(y));
    }
    ov_label:
        ...
}
```

- We mentioned spaghetti code...

Hardware Support

- Motivating and rejected examples
- Survey of software interfaces
- **Hardware support**
 - Existing hardware: flags
 - Existing hardware: traps
 - Flags versus traps
- Hardware / software mapping
- A scarecrow proposal

Existing HW: Flags

- Basic operations:
 - Save registers
 - Restore registers
 - Test flags
- One or more registers visible in ISA
 - May include “last instruction” flags
- May be additional internal storage
 - e.g. with reorder buffer entry

Existing HW: Traps

- Basic operations:
 - Enable trapping
 - Disable trapping
 - Set handler
- Currently require OS support
 - Need privileged mode to set handler
 - Handler runs in privileged mode
- Trap enable/disable on IA32 costs more than flag save/restore

Flags versus Traps

- Traps are an optimization for flag test and branch
 - But flag tests are reasonably inexpensive!
 - Flag tests need only occur at synchronization points (identified by programmer or compiler)
- There are other possible optimizations:
 - Execution predicated on flag settings
 - Conditional branch FP ops
 - And others...
- Compiler could optimize away explicit tests

HW/SW Mapping

- Motivating and rejected examples
- Survey of software interfaces
- Hardware support
- **Hardware / software mapping**
 - Extended range: a case study
 - fpmenu implementation notes
 - Interfaces and implementations
 - Performance
- A scarecrow proposal

Extended Range: A Case Study

```
scaled_double prod;  
for (i = 0; i < n; ++i)  
    prod *= a[i];
```

Extend range by implementing a scaled precision:

- No exceptions: scale on every operation
- Flags: test after each operation
- Traps: use “counting mode”

Can optimize first two cases by blocking.

Extended Range: A Case Study

```
scaled_double prod;
for (i = 0; i < n; i += BLOCK) {
    prod_tmp = fast product over block
    if (no range exception)
        prod *= prod_tmp;
    else
        prod *= scaled subproduct
}
```

- Compiler ideally generates this from previous code
- Otherwise, little worse than blocking matrix codes
 - Could probably use similar automatic tuning

HW/SW Mapping

fpmenu:

- Uses SIGFPE handler + ugly C macros to implement try/catch and replacement
- On exception
 - try/catch: restore state, jump to user
 - substitution: decode, compute, writeback
 - other: re-execute instruction

HW/SW Mapping

fpmenu:

- handler choice really needs compiler input
- must manually add fwait instructions
- makes most optimizations dangerous
- context save penalty on try/catch entry
- instruction re-execution and toggling traps are both expensive

Compiler support would help, but some problems are intrinsic to trap-based handling.

Interfaces and Implementations

Several implementations for software interfaces

- Compile flag test and branch to trapping code
- Implement try-catch handling with flag tests

Software resources need not map directly to HW

- Map HW invalid flag to multiple software flags
- Support flag-carrying types with virtualized flags
- Merge local HW flags into virtual global register

We standardize *interface* requirements, not implementations. Simple basic interfaces are easier to reason about and permit adequate room to optimize.

Performance: Traps v. Flags

- Platforms tested: PPro@233MHz, P3@800Mhz, and P4@1.4GHz
- Tested continued fractions, long products.
- Results: Blocked flag tests are fine.
 - Blocked flag tests usually faster than trapping.
 - Immediate flag tests are considerably slower.
 - Detrimental in tight loops like long product.
 - Same as trapping in continued fractions.
 - Saving machine state for try-catch is slow.

Scarecrow Proposal

- Motivating and rejected examples
- Survey of software interfaces
- Hardware support
- Hardware / software mapping
- **A scarecrow proposal**

Scarecrow Proposal

A scarecrow is a spooky outline of a straw-man.

SHALL be able to tie flags to value types.

- Compilers already need this for optimizations.
- Can use this to deliver information on exceptions affecting results.
- Generalizes to vectors in many ways.

Scarecrow Proposal

Helpers:

SHALL provide a presubstitution mechanism

- Presubstitute the result of an *expression*.
- Any unhandled conditions remain raised.
- Types are known.

SHALL provide scaled-exponent type

- Many uses, can be heavily optimized.
- Must explicitly determine the substitute's sign.
- Note that doubled-precisions don't round correctly.

Scarecrow Proposal

SHALL require programmers to declare interest in flags

- Scope and extent?
- All flags or particular flags?

SHALL provide an invalid hierarchy

MAY allow users to add conditions

-
- Borneo's admit-yields is a good example.
 - Static scope and extent flag declarations make user-defined, software flags reasonable.

Scarecrow Proposal

edit Move exception handling optimizations to an informative annex.

- Describe fast trapping or conditional operations as optimizations

edit Eliminate signaling NaNs.

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- Only current use of signaling NaNs: debugging.

End.

You're still here? Go home.