

*Sparse Data Structures for Weighted Bipartite  
Matching*

E. Jason Riedy  
Dr. James Demmel  
(... and thanks to the BeBOP group)

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## *Use Sparse Matrix Optimizations...*

- ▶ Take a fixed, simple algorithm: Auction alg. for matchings
  - ▶ Repeated iterations over a sparse graph.
- ▶ What's expensive, and is there anything we can do about it?
  - ▶ Take an idea from optimizing sparse matrix-vector products.
- ▶ A little speed-up in some cases, but there are more ideas available...

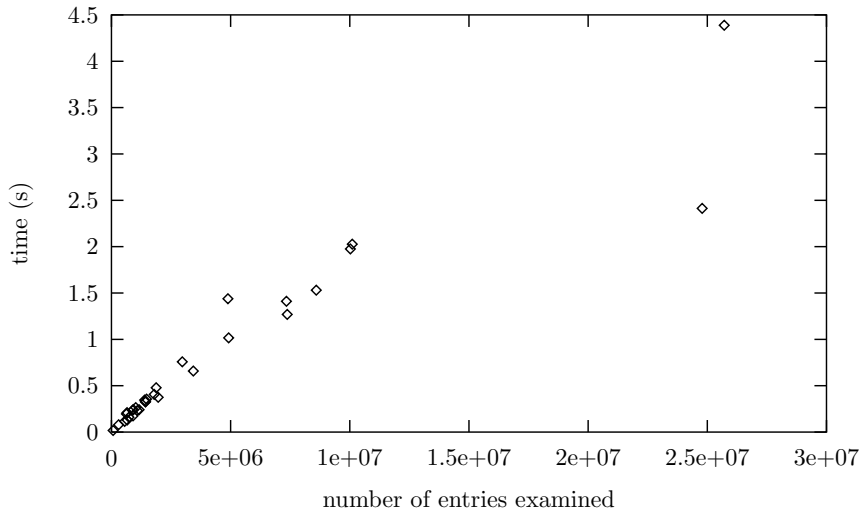
## *Where's the Time Going?*

Auction algorithm: Iterative, greedy algorithm bipartite matching:

1. An unmatched row  $i$  finds a “most profitable” column  $j$ 
  - ▶  $\pi(i) = \max_j b(i, j) - p(i)$
2. Row  $i$  places a bid for column  $j$ .
  - ▶ Bid price raised until  $j$  is no longer the best choice. (Min. increment  $\mu$ )
3. Highest bid gets the matching  $(i, j)$ .

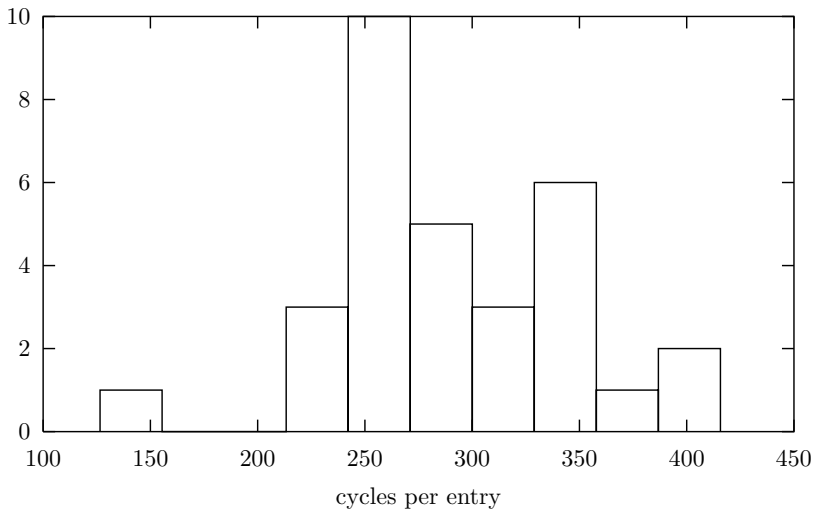
## *Time linear in entries examined...*

Number of entries examined is problem-dependent.



## *Expensive Inner Loop!*

1.3 GHz Itanium 2








## *And Locating...*

No obvious culprits in the instructions...

5	<input type="text" value="4.62"/>	4.62	int j = ind[k];
6	<input type="text" value="4.62"/>	4.62	double e = ent[k];
7			double p = price[j];
3			
3	<input type="text" value="13.87"/>	13.87	val = e - p;
2			
1	<input type="text" value="9.25"/>	9.25	if (val <= best_val_1) continue;

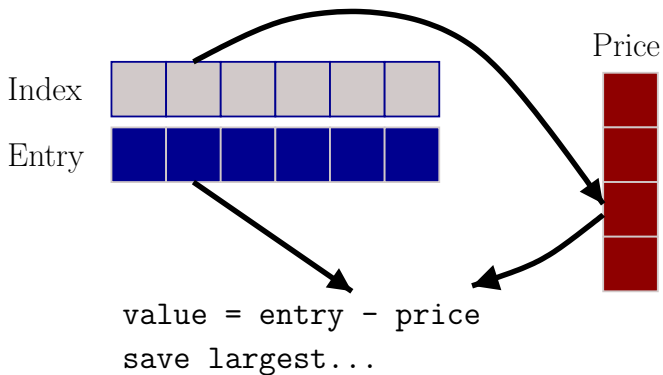
## *And Locating...*

But considering cache effects!

<pre>5</pre>		31.81	<pre>int j = ind[k];</pre>
<pre>6</pre>		51.57	<pre>double e = ent[k];</pre>
<pre>7</pre>			<pre>double p = price[j];</pre>
<pre>8</pre>			
<pre>9</pre>			
<pre>10</pre>		16.62	<pre>val = e - p;</pre>
<pre>11</pre>			

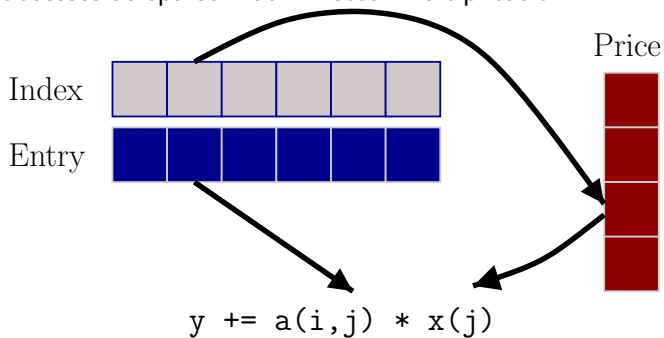


## *Auction's Inner Loop*

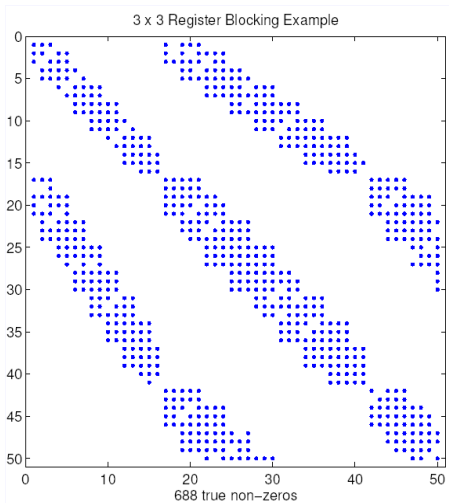


## *Auction's Inner Loop*

Same accesses as sparse matrix-vector multiplication!

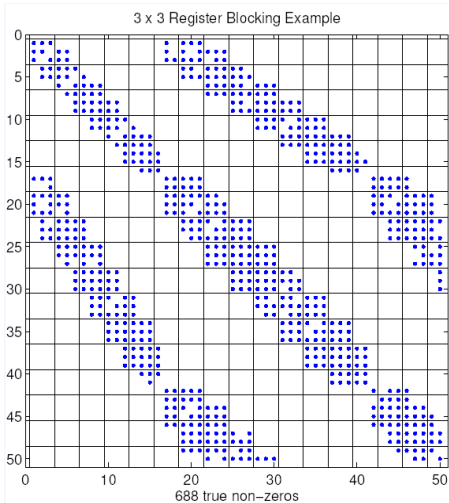


## *Performance Through Blocking?*



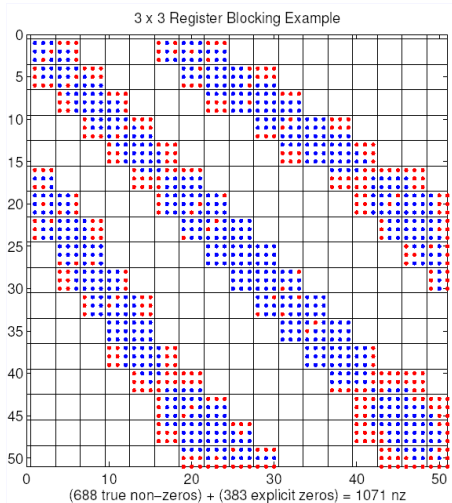
(Images swiped from Berkeley's BeBOP group.)

## Performance Through Blocking?



(Images swiped from Berkeley's BeBOP group.)

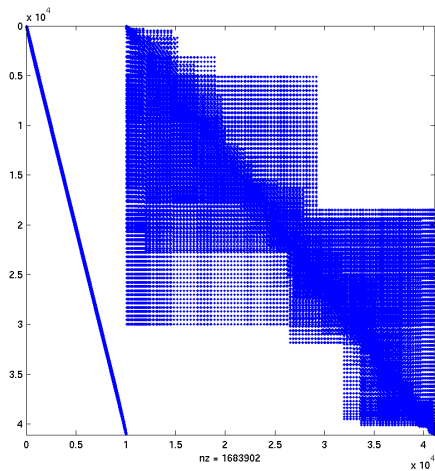
## Performance Through Blocking?



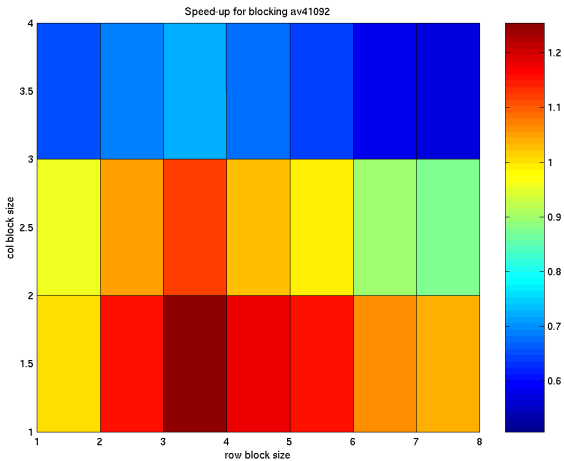
More entries, but  $1.5\times$  performance on Pentium 3!  
(Images swiped from Berkeley's BeBOP group.)

## *Blocking Speeds Some Matches*

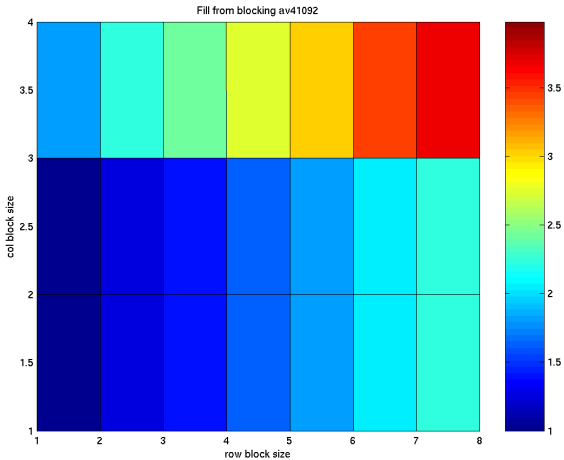
Finite element matrix from Vavasis (in UF collection):



# *Blocking Speeds Some Matches*



# *Blocking Speeds Some Matches*





## Observations

A blocked graph data structure may provide additional performance if:

- ▶ you iterate over whole rows,
- ▶ the graph / matrix has runs of columns, and
- ▶ you're willing to use an automated tuning system.

Maximizing the runs: linear arrangement. Hard, but there may be cheap heuristics. Only worth-while if you're performing **many** iterations. (For mat-vec, often  $> 50$  computations of  $Ax$ .)